Mathematical Aspects of Mixing Times in Markov Chains

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Mathematical Aspects of Mixing Times in Markov Chains

Ravi Montenegro¹ and Prasad Tetali²

Abstract

In the past few years we have seen a surge in the theory of finite Markov chains, by way of new techniques to bounding the convergence to stationarity. This includes functional techniques such as logarithmic Sobolev and Nash inequalities, refined spectral and entropy techniques, and isoperimetric techniques such as the average and blocking conductance and the evolving set methodology. We attempt to give a more or less self-contained treatment of some of these modern techniques, after reviewing several preliminaries. We also review classical and modern lower bounds on mixing times. There have been other important contributions to this theory such as variants on coupling techniques and decomposition methods, which are not included here; our choice was to keep the analytical methods as the theme of this presentation. We illustrate the strength of the main techniques by way of simple examples, a recent result on the Pollard Rho random walk to compute the discrete logarithm, as well as with an improved analysis of the Thorp shuffle.

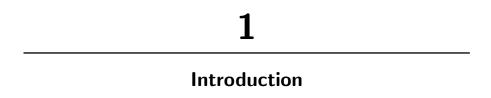
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Monte Carlo methods have been in use for a long time in statistical physics and other fields for sampling purposes. However, the computer scientists' novel idea [38] of reducing the problem of approximately counting the size of a large set of combinatorial objects to that of near-uniform sampling from the same set, gave the study of Markov chain Monte Carlo (MCMC) algorithms an entirely new purpose, and promptly spawned an active subtopic of research. We recall here that the work of [38] shows that in fact, under the technical assumption of so-called *self-reducibility*, approximate counting of the size of a set in *polynomial time* is feasible if and only if one is able to sample from the set with nearly uniform distribution, also in polynomial time. In terms of the finite Markov chain underlying an MCMC algorithm, the latter problem translates to designing and analyzing a Markov chain with a prescribed stationary measure, with a view (and hope) to providing rigorous estimates on the polynomial fastness of the rate of convergence to stationarity of the chain. Thus the classical subject of finite Markov chains has received much renewed interest and attention.

To add concreteness to the above story, we briefly mention as examples of large sets of combinatorial objects, the set of matchings of a given

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(as input) bipartite graph [34, 36], the set of proper colorings of a given graph using a fixed number of colors [28], the number of matrices having non-negative integer entries and with prescribed row and column sums [16], etc. Albeit combinatorial, a non-discrete estimation problem which received significant devotion, both by way of algorithms and analytical techniques, is that of (approximately) computing the volume of a high-dimensional convex body (see [47, 48] and references therein). There have already been some very good surveys focusing on such combinatorial, computational and statistical physics applications of finite Markov chains. For an elaboration of the above premise, and a crash course on several basic techniques, we recommend the excellent article of Jerrum [32]. Towards the end of this introduction, we provide other pointers to existing literature on this subject. However, much of the theory surveyed here is rather recent theoretical (analytical) development and is so far unavailable in a unified presentation. The significance of these new methods is as follows.

The rate of convergence to stationarity of a finite Markov chain is typically measured by the so-called mixing time, defined as the first time τ by which the L^1 (or more generally, L^p) distance between the distribution at time τ and the stationary distribution falls below a small threshold, such as 1/2e. It is classical and elementary to show that the inverse spectral gap of a lazy reversible Markov chain captures the mixing time (in L^1 and L^2) up to a factor of $\log(1/\pi_*)$, where $\pi_* = \min_x \pi(x)$ denotes the smallest entry in the stationary probability (vector) π of the chain. While the more technical logarithmic Sobolev constant captures the L^2 -mixing time up to a factor of $\log \log(1/\pi_*)$, it is typically much harder to bound – to mention a specific example, the exact constant is open for the 3-point space with arbitrary invariant measure; also in a few cases, the log-Sobolev constant is known not to give tight bounds on the L^1 -mixing time. The main strength of the spectral profile techniques and the evolving set methodology considered in this survey seems to be that of avoiding extra penalty factors such as $\log \log(1/\pi_*)$. These extra pesky factors can indeed be non-negligible when the state space is of exponential (or worse) size in the size of the input. In the present volume, the above is illustrated with a couple of simple examples, and with the now-famous Thorp shuffle, for which an improved $O(d^{29})$ mixing time is described, building on the proof of Morris that proved the first polynomial (in d) bound of $O(d^{44})$ – here the number of cards in the deck is 2^d , and hence the state space has size $2^d!$, resulting in a $\log \log(1/\pi_*)$ factor of only O(d), while a $\log(1/\pi_*)$ factor would have yielded an all too costly $O(d2^d)$.

The approach to L^2 -mixing time using the spectral profile has the additional advantage of yielding known (upper) estimates on mixing time, under a log-Sobolev inequality and/or a Nash-type inequality. Thus various functional analytic approaches to mixing times can be unified with the approach of bounding the spectral profile. The one exception to this is the approach to stationarity using relative entropy; the corresponding *entropy constant* capturing the rate of decay of entropy has also been hard to estimate.

A brief history of the above development can perhaps be summarized as follows. A fundamental contribution, by way of initiating several subsequent works, was made by Lovász and Kannan in [46] in which they introduced the notion of *average conductance* to bound the total variation mixing time. This result was further strengthened and developed by Morris and Peres using the so-called *evolving sets*, where they analyze a given chain by relating it to an auxiliary (dual) chain on subsets of the states of the original chain. While this was introduced in [61] in a (martingale-based) probabilistic language, it turns out to be, retrospectively, an independent and alternative view of the notion of a Doob transform introduced and investigated by Diaconis and Fill [18]. Further refinement and generalization of the evolving sets approach was done in detail by [56]. The functional analog of some of this is done via the spectral profile, developed for the present context of finite Markov chains, in [30], while having its origins in the developments by [4] and [14] in the context of manifolds.

Besides summarizing much of the above recent developments in this exciting topic, we address some classical aspects as well. In discretetime, much of the literature uses laziness assumptions to avoid annoying technical difficulties. While laziness is a convenient assumption, it slows down the chain by a factor of 2, which may not be desirable in practice. We take a closer look at this issue and report bounds which reflect the precise dependence on laziness. The notion of modified conductance circumvents laziness altogether, and we discuss this aspect briefly and compare it to bounds derived from the functional approach.

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Further details on the modified conductance and its usefulness can be found in [57]. Another issue is that of the role of reversibility (a.k.a. detailed balance conditions). We tried to pay particular attention to it, due to current trend in the direction of analyzing various nonreversible Markov chains. Although often a convenient assumption, we avoid as much as possible this additional assumption. In particular, we include a proof of the lower bound on the total variation mixing time in terms of the second eigenvalue in the general case. Besides providing upper and lower bounds for the mixing time of reversible and non-reversible chains, we report recent successes (with brief analysis) in the analysis of some non-reversible chains; see for example, the Pollard Rho random walk for the discrete logarithm problem and the Thorp shuffle.

In Section 2 we introduce notions of mixing times and prove the basic upper bounds on these notions using Poincaré and logarithmic Sobolev type functional constants. In Section 3 we move on to recent results using the spectral profile, as opposed to using simply the second eigenvalue. In Section 4 we review the evolving set methods. Our treatment of lower bounds on mixing times is provided in Section 5. We consider several examples for illustration in Section 6. In the penultimate section, we gather a few recent results together. This includes recent results on the so-called fastest mixing Markov chain problem, and a recent theorem [52] from perturbation theory of finite Markov chains; this theorem relates the stability of a stochastic matrix (subject to perturbations) to the rate of convergence to equilibrium of the matrix. We also recall here an old but not so widely known characterization of the spectral gap, which seems worth revisiting due to recent results utilizing this formulation. The Appendix contains a discussion on the relations between the distances considered in this survey, and others such as relative pointwise (L^{∞}) distance.

We mention here a few additional sources, by way of survey articles, for the interested reader. For a good overview of the basic techniques in estimating the mixing times of finite Markov chains, see [32, 33, 35]. Other updates include the tutorial lectures of [40, 65]. Also a recent manuscript of Dyer et al. [25] describes several comparison theorems for reversible as well as nonreversible Markov chains.

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Appendix

Our main focus has been on bounds for L^2 distance. Our bounds on mixing times in L^2 and relative entropy also yield bounds on the total variation mixing time using the following well-known inequality relating probability measures ν and μ .

$$\|\nu - \mu\|_{\rm TV} = \frac{1}{2} \left\|\frac{\nu}{\mu} - 1\right\|_{1,\mu} \le \frac{1}{2} \left\|\frac{\nu}{\mu} - 1\right\|_{2,\mu}.$$
 (8.1)

Further assuming that ν is absolutely continuous with respect to μ , the so-called Pinsker inequality (see Lemma 12.6.1 in [15] for a proof), asserts that:

$$\|\nu - \mu\|_{\rm TV}^2 \le \frac{1}{2} D(\nu \|\mu) \tag{8.2}$$

Finally the general inequality $(\mathbb{E}_{\mu}f)\operatorname{Ent}_{\mu}(f) \leq \operatorname{Var}_{\mu}(f)$, valid for all measurable functions on an arbitrary probability space (since $\log \frac{f}{\mathbb{E}_{\mu}f} \leq \frac{f}{\mathbb{E}_{\mu}f} - 1$), when applied to $f = \nu/\mu$ implies that,

$$D(\nu \| \mu) \le \left\| \frac{\nu}{\mu} - 1 \right\|_{2,\mu}^2.$$
 (8.3)

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In a sense the L^{∞} , or relative pointwise distance, is the strongest of all distances. Our L^2 bounds also induce L^{∞} bounds. Observe that if $t = t_1 + t_2$ then

$$\left|\frac{H_t(x,y) - \pi(y)}{\pi(y)}\right| = \left|\frac{\sum_z (H_{t_1}(x,z) - \pi(z))(H_{t_2}(z,y) - \pi(y))}{\pi(y)}\right|$$
$$= \left|\sum_z \pi(z) \left(\frac{H_{t_1}(x,z)}{\pi(z)} - 1\right) \left(\frac{H_{t_2}^*(y,z)}{\pi(z)} - 1\right)\right|$$
$$\leq \left\|h_{t_1}^x - 1\right\|_2 \left\|h_{t_2}^{*y} - 1\right\|_2$$
(8.4)

where the inequality follows from Cauchy-Schwartz. Several bounds on L_{∞} mixing then follow immediately, including

$$au_2(\epsilon) \le au_\infty(\epsilon) \le au_2\left(\epsilon\sqrt{\frac{\pi_*}{1-\pi_*}}\right)$$

and

$$\tau_2(\epsilon) \le \tau_{\infty}(\epsilon) \le \tau_2^{\mathsf{P}}(\sqrt{\epsilon}) + \tau_2^{\mathsf{P}^*}(\sqrt{\epsilon})$$

The first of these two is somewhat unappealing because the asymptotic portion of $\tau_2(\epsilon)$ is of the form $\lambda^{-1}\log(1/\epsilon)$, and so taking $\tau_2\left(\epsilon\sqrt{\frac{\pi_*}{1-\pi_*}}\right)$ adds an extra factor of $\lambda^{-1}\log(1/\pi_*)$ to the $\tau_2(\epsilon)$ bound, potentially large relative to spectral profile bounds. The second bound unfortunately requires study of both P and P*. However, if P is reversible then this last bound becomes

$$\tau_2(\epsilon) \le \tau_\infty(\epsilon) \le 2\tau_2(\sqrt{\epsilon}). \tag{8.5}$$

More generally, most bounds in this survey were the same for P and P^{*}. For instance, (8.5) holds for the spectral profile bounds on L^2 mixing in terms of $\Lambda(r)$. In particular, the Dirichlet form satisfies

$$\mathcal{E}_{\mathsf{P}}(f,f) = \mathcal{E}_{\mathsf{P}^*}(f,f)$$

and so $\lambda(\mathsf{P}) = \lambda(\mathsf{P}^*)$, $\Lambda_{\mathsf{P}}(r) = \Lambda_{\mathsf{P}^*}(r)$, $\rho(\mathsf{P}) = \rho(\mathsf{P}^*)$ and $\Phi_{\mathsf{P}}(r) = \Phi_{\mathsf{P}^*}(r)$ (as $\mathsf{Q}(A, A^c) = \mathcal{E}(\mathbf{1}_A, \mathbf{1}_A)$).

It is not as clear how the *f*-congestion bounds behave for P^* . However, if π is uniform then

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$$\Psi(A) = \Psi(A^c) = \min_{\pi(B)=\pi(A)} \mathsf{Q}(A^c, B)$$
$$= \min_{\pi(B)=\pi(A)} \mathsf{Q}_{\mathsf{P}^*}(B, A^c) \ge \min_{\pi(B)=\pi(A)} \Psi_{\mathsf{P}^*}(B)$$

and so $\tilde{\phi}(r) \geq \tilde{\phi}_{\mathsf{P}^*}(r)$. The converse follows similarly, so $\tilde{\phi}_{\mathsf{P}}(r) = \tilde{\phi}_{\mathsf{P}^*}(r)$ when π is uniform, and (8.5) holds for modified-conductance bounds.